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*One Combinatorial Algorithm  
for Berth Allocation Problem*

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# Input Data

**$T$** : Total number of time periods in the planning horizon.

**$m$**  : The number of berths in the port.

**$l$**  : The number of vessels in the planning horizon.

***vessel***: Sequence of data relevant for vessels, it has following structure:

$$vessel = \{ (ETA_k, a_k, b_k, d_k, s_k, c_{1k}, c_{2k}, c_{3k}, c_{4k}) \mid k = 1 \dots l \}.$$

Members of the 9-tuple represents following data of the vessel:

$ETA_k$ : Expected time of arrival of  $vessel_k$ ;

$a_k$ : The processing time of  $vessel_k$  if only one crane is assigned to  $vessel_k$ ;

$b_k$ : The length of  $vessel_k$ ;

$d_k$ : The due time for the departure of  $vessel_k$ ;

$s_k$ : The least-cost berthing location of the reference point of  $vessel_k$ ;

$c_{1k}$ : The penalty cost of  $vessel_k$  if the vessel could not dock at its preferred berth;

$c_{2k}$ : The penalty cost of  $vessel_k$  per unit time of earlier arrival before  $ETA_k$ ;

$c_{3k}$ : The penalty cost of  $vessel_k$  per unit time of late arrival after  $ETA_k$ ;

$c_{4k}$ : The penalty cost of  $vessel_k$  per unit time delay behind the due time  $d_k$ .

### 3.1.2 Decision variables and domains

$At_k$  : The arrival time of vessel  $k$  to the berth.

Domain ( $At_k$ )= $\{1,2,3,4,\dots,T\}$

$Dt_k$  : The departing time of vessel  $k$ .

Domain ( $Dt_k$ )= $\{1,2,3,4,\dots,T\}$

$X_{itk}$  : 1 if the berth  $i$  at time  $t$  is allocated to vessel  $k$ , otherwise 0.

Domain ( $X_{itk}$ )= $\{0,1\}$

### 3.1.3 Constraints

**Constraint 3-1-1:** The grid squares are covered by only one vessel. In fact, each berth at time t can be assigned to only one vessel.

$$\sum_{k=1}^l X_{itk} \leq 1 \text{ for } i = 1, 2, 3, \dots, m; t = 1, 2, 3, \dots, T$$

**Constraint 3-1-2:** Each berth is allocated for the vessel only between its arrival and departure.

$$\begin{aligned} At_k \leq t \leq Dt_k &\Rightarrow X_{itk} = 1 \\ (At_k > t) \text{ OR } (t > Dt_k) &\Rightarrow X_{itk} = 0 \\ \text{for } t = 1, 2, \dots, T; \text{ for } i = 1, 2, \dots, m; \text{ for } k = 1, 2, \dots, l \end{aligned}$$

### 3.1.4 Objective function

The objective function of this decision is to minimize the total penalty costs. In order to present the objective function, we introduce the following auxiliary variable:

$Z_k$  : The sum of the absolute distance between the preferred location of vessel  $k$  and the berths allocated to the vessel. This variable is determined by the following equation:

$$Z_k = f(X_{itk}, s_k) = \sum_{t=1}^T \sum_{i=1}^m \{ |i - s_k| : X_{itk} = 1 \}$$

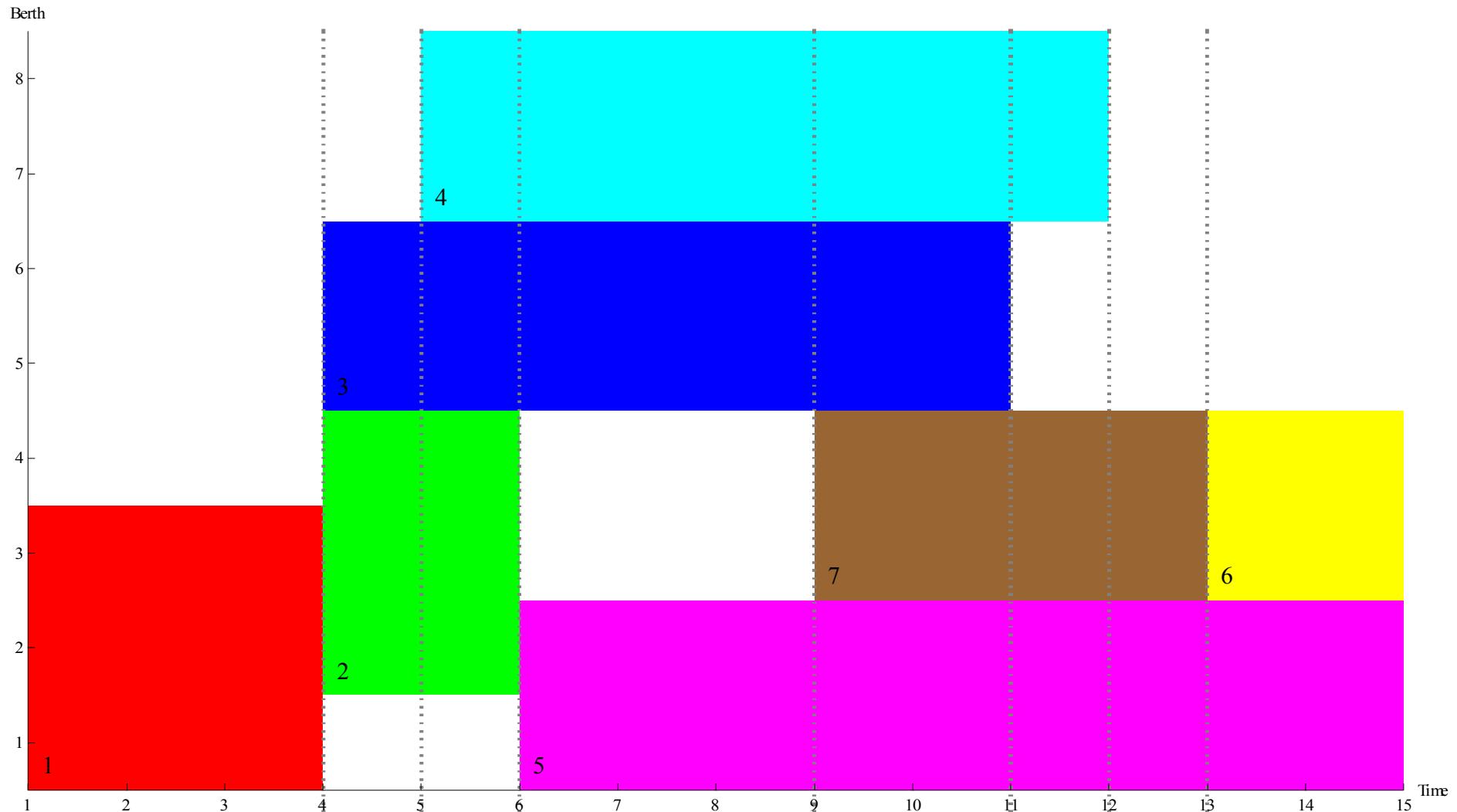
Now the objective function is written as follows:

$$\text{MinCostVessels} = \sum_{k=1}^l \{ c_{1k} \cdot Z_k + c_{2k}(ETA_k - At_k)^+ + c_{3k}(At_k - ETA_k)^+ + c_{4k}(Dt_k - d_k)^+ \}$$

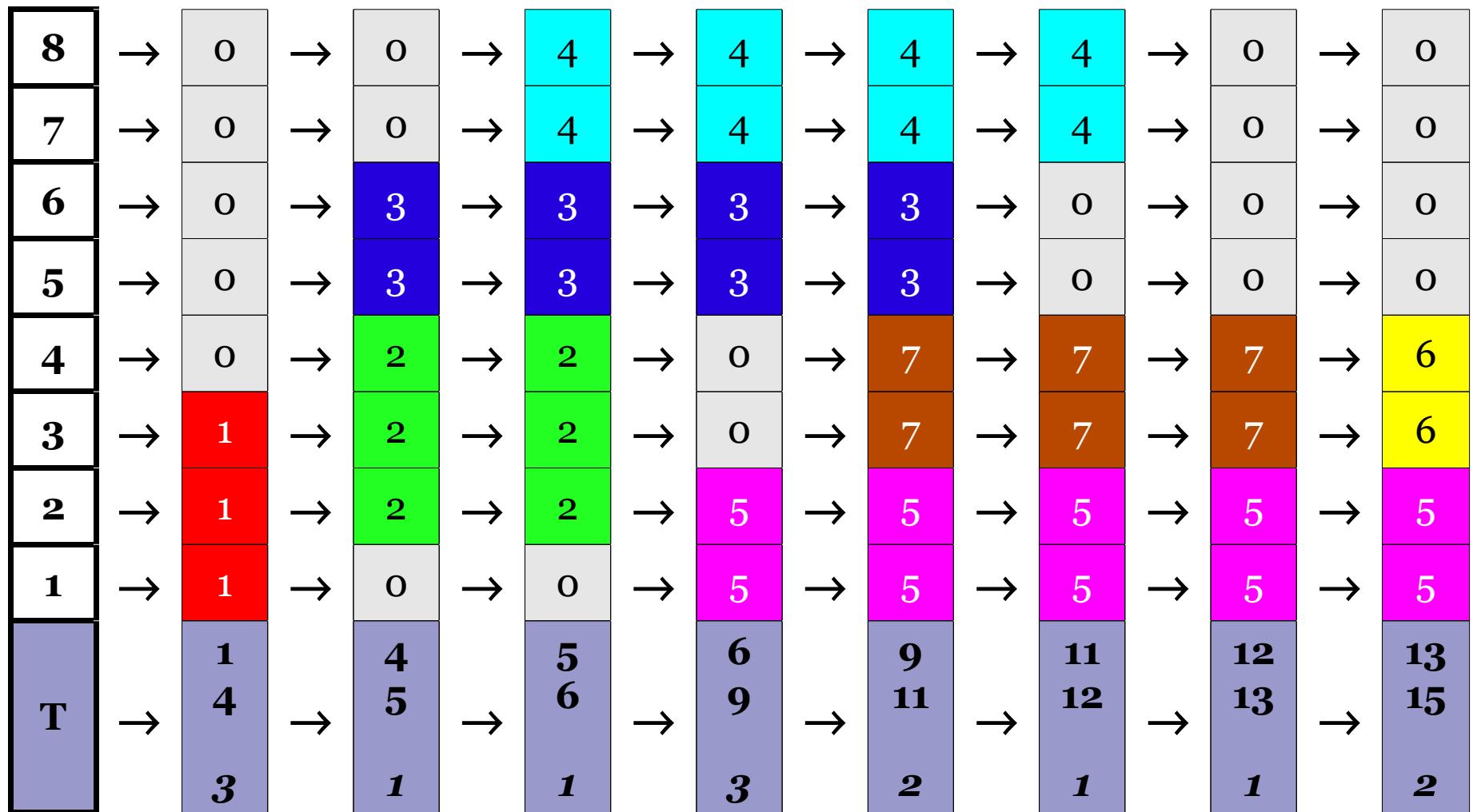
## Example of input data

$k$	$ETA_k$	$a_k$	$b_k$	$d_k$	$s_k$	$c_{1k}$	$c_{2k}$	$c_{3k}$	$c_{4k}$
<b>1</b>	1	9	3	4	1	25	10	10	30
<b>2</b>	4	6	3	6	3	40	30	20	35
<b>3</b>	4	14	2	11	5	10	20	20	25
<b>4</b>	5	14	2	12	7	15	20	20	30
<b>5</b>	6	18	2	16	2	35	25	20	45
<b>6</b>	13	4	2	16	4	20	30	35	40
<b>7</b>	10	8	2	14	3	25	10	10	30
<b>8</b>	6	6	2	9	4	40	35	40	70
<b>9</b>	2	8	4	2	2	20	30	30	25

# Example of output data



# Internal list data structure representation



# Distribution of cost for a vessel with ETA=3 at berth 5

<b>m</b>							...		
<b>m-1</b>							...		
.	.	.	.	.	.	.	.	.	.
.	.	.	.	.	.	.	.	.	.
.	.	.	.	.	.	.	.	.	.
<b>7</b>	50	40	30	40	50	60	...		
<b>6</b>	40	30	20	30	40	50	...		
<b>5</b>	30	20	<b>10</b>	20	30	40	...		
<b>4</b>	40	30	20	30	40	50	...		
<b>3</b>	50	40	30	40	50	60	...		
<b>2</b>	60	50	40	50	60	70	...		
<b>1</b>	70	60	50	60	70	80	...		
Berth Time	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>	<b>6</b>	...	<b>T-1</b>	<b>T</b>

## Sequence $\xi$ for a vessel with ETA=4 at berth 5

$$\begin{aligned}\xi = \{ & (5,3,10), \\ & (4,3,20), (5,2,20), (5,4,20), (6,3,20), \\ & (3,3,30), (4,2,30), (4,4,30), \dots (7,3,30), \\ & \cdot \\ & \cdot \\ & \cdot \\ & (m,T,x) \} \end{aligned}$$

## Brute Force Algoritam

```
MinCost = +∞;
n = Table[0, {ID,1}];
ID = 1;
λ = 1;
stack = ξ;
While 0 < ID Do
    DeleteVessel[ID];
    If λ == -1 Then
        n[ID] = 0; ID--;
        Pop[];
        λ = 1,
    Else
        If n[ID] < Length[ξ[ID]] Then
            n[ID]++;
            berth = ξ[ID,n[ID],1];
            arrival = ξ[ID,n[ID],2];
            tempCost = Sum[ξ[i,n[i],3],{i,ID}];
            InsertVessel[ID,berth,arrival];
            If ID == 1 Then
                ReportSolution[tempCost];
                λ = -1
            Else
                If tempCost < minCost Then
                    Push[];
                    ID++
                Else
                    λ = -1
                EndIf
            EndIf
        Else
            λ = -1
        EndIf
    EndIf
EndWhile
```

# *Algoritam Opt 1*

```
MinCost = +∞;
n = Table[0, {ID,1}];
ID = 1;
λ = 1;
stack = ξ;
While 0 < ID Do
    DeleteVessel[ID];
    If λ == -1 Then
        n[ID] = 0; ID--;
        Pop[];
        λ = 1,
    Else
        If n[ID] < Length[ξ[ID]] Then
            n[ID]++;
            berth = ξ[ID,n[ID],1];
            arrival = ξ[ID,n[ID],2];
            tempCost = Sum[ξ[i,n[i],3],{i,ID}];
            InsertVessel[ID,berth,arrival];
            If ID == 1 Then
                ReportSolution[tempCost];
                λ = -1
            Else
                If tempCost + minVesselsCosts[] < minCost Then
                    Push[];
                    If tempCost + minVesselsCosts[] < minCost Then ID++
                    Else Pop[]
                EndIf
                Else
                    λ = -1
                EndIf
            EndIf
        Else
            λ = -1
        EndIf
    EndIf
EndWhile
```

# Configuration of the principal variables of the *Algorithm Opt 1*

ID	1	2	3	4	5	l-1	l	
$\xi$			(8,11,115)	(7,10,96)				
		(1,9,100)	(8,11,115)	(6,10,96)	(8,11,115)		(8,11,115)	
	(8,9,100)	(2,9,95)	(7,9,110)	(7,9,96)	(7,9,110)	.	(6,9,110)	
	(8,9,100)	(1,8,85)	(7,9,110)	(6,9,96)	(7,9,110)	.	(6,8,110)	
	.	.	.	.	.	.	.	
	.	.	.	.	.	.	.	
	.	.	.	.	.	.	.	
	(3,1,30)	(6,3,15)	(5,4,20)	(5,1,8)	(7,7,6)	.	(6,4,10)	
	(1,2,20)	(7,4,15)	(5,2,20)	(6,2,8)	(8,8,6)	.	(7,2,10)	
	(2,1,20)	(5,4,15)	(4,3,20)	(5,3,8)	(6,8,6)	.	(5,3,10)	
n	(1,1,10)	(6,4,5)	(5,3,10)	(4,3,4)	(7,8,3)	.	(6,3,10)	
	2	3	0	0	0	.		
	20	15	10	4	3	.	8	
$\sum$	TempCost = 35		minVesselsCost[]					

*Step Forward* – Selection of the position (5,3,10) for the ID=3 vessel

*Step Forward – Deletion* of the position (5,3,10) for vessels with ID>3

*Step Forward – Adjustment of the sums. Ready for the next step.*

ID	1	2	3	4	5	l-1	l
$\xi$			(8,11,115)				
		(1,9,100)	(8,11,115)	(7,10,96)	(8,11,115)		
	(8,9,100)	(2,9,95)	(7,9,110)	(6,10,96)	(7,9,110)	.	.
	(8,9,100)	(1,8,85)	(7,9,110)	(7,9,96)	(7,9,110)	.	.
	.	.	.	.	.	.	.
	.	.	.	.	.	.	.
	.	.	.	.	.	.	.
	(3,1,30)	(6,3,15)	(5,4,20)	(4,3,8)	(7,7,6)	.	.
	(1,2,20)	(7,4,15)	(5,2,20)	(5,1,8)	(8,8,6)	.	.
	(2,1,20)	(5,4,15)	(4,3,20)	(6,2,8)	(6,8,6)	.	.
n	(1,1,10)	(6,4,5)	(5,3,10)	(4,3,4)	(7,8,3)	.	.
	2	3	1	0	0	.	.
	20	15	10	4	3	.	.
$\sum$	TempCost = 45			minVesselsCost[]			

# T=15, m=8 – Brute Force Algorithm vs. Algorithm Opt 1

$l$	Brute Force Algorithm		Algorithm Opt 1		$\times$
	Time	No. Sol.	Time	No. Sol.	
6	1.341	3	0.094	2	<b>14.266</b>
7	7.550	11	0.546	8	<b>13.828</b>
8	13.712	25	0.702	7	<b>19.533</b>
9	19.687	25	1.076	10	<b>18.296</b>
10	128.981	30	12.823	6	<b>10.058</b>
11	357.340	224	45.458	10	<b>7.861</b>
12	472.728	54	59.982	11	<b>7.881</b>
13	651.644	345	100.448	18	<b>6.487</b>
14	847.050	550	130.556	16	<b>6.488</b>
15	2245.228	853	276.713	27	<b>8.114</b>

## Algoritam Opt 2 × 2

```
MinCost = +∞;
n = Table[0, {ID,1}];
ID = 1;
λ = 1;
stack = ξ;
While 0 < ID Do
    DeleteVessel[ID];
    If λ == -1 Then
        n[ID] = 0; ID--;
        Pop[];
        λ = 1,
    Else
        If n[ID] < Length[ξ[ID]] Then
            n[ID]++;
            berth = ξ[ID,n[ID],1];
            arrival = ξ[ID,n[ID],2];
            tempCost = Sum[ξ[i,n[i],3],{i,ID}];
            InsertVessel[ID,berth,arrival];
            If ID == 1 Then
                ReportSolution[tempCost];
                λ = -1
            Else
                If tempCost + minVesselsCosts[] < minCost Then
                    Push[];
                    Heuristics[];
                    If tempCost + minVesselsCosts[] < minCost Then ID++
                    Else Pop[]
                EndIf
                Else
                    λ = -1
                EndIf
            EndIf
        Else
            λ = -1
        EndIf
    EndIf
EndWhile
```

# **T=15, m=8 – Algorithm Opt 1 vs. Algorithm Opt 2 $\times$ 2**

$l$	Algorithm Opt 1		Algorithm Opt 2 $\times$ 2		$\times$	$\times_{BF}$
	Time	No. Sol.	Time	No. Sol.		
6	0.094	2	0.094	2	<b>1.000</b>	<b>14.265</b>
7	0.546	8	0.234	3	<b>2.333</b>	<b>32.265</b>
8	0.702	7	0.437	6	<b>1.606</b>	<b>31.378</b>
9	1.076	10	1.326	8	<b>0.811</b>	<b>14.846</b>
10	12.823	6	6.614	5	<b>1.939</b>	<b>19.501</b>
11	45.458	10	16.380	9	<b>2.775</b>	<b>21.816</b>
12	59.982	11	22.354	8	<b>2.683</b>	<b>21.147</b>
13	100.448	18	35.490	10	<b>2.830</b>	<b>18.361</b>
14	130.556	16	47.221	10	<b>2.765</b>	<b>17.938</b>
15	276.713	27	111.856	13	<b>2.474</b>	<b>20.073</b>

## *CPLEX 11.2*

Linux Slackware 12, Kernel: 2.6.21.5

Intel Core 2 Duo CPU E6750 on 2.66GHz

RAM=8Gb

## *Algorithm Opt 2×2*

Microsoft Windows 7 64bit OS

Wolfram Mathematica 8.0 (64 bit version)

HP Pavilion dv3 Computer

Intel Core i3 CPU M350 @2.27GHz

RAM=4Gb.

# **T=15, m=8 – CPLEX 11.2 vs. Algorithm Opt $2 \times 2$**

$l$	CPLEX 11.2	Algorithm Opt $2 \times 2$		$\times$
	Time	Time	No. Sol.	
6	0.06	0.094	2	<b>0.638</b>
7	20.53	0.234	3	<b>87.735</b>
8	18.91	0.437	6	<b>43.272</b>
9	20.88	1.326	8	<b>15.747</b>
10	35.19	6.614	5	<b>5.321</b>
11		16.380	9	
12	129.76	22.354	8	<b>5.805</b>
13	379.64	35.490	10	<b>10.697</b>
14		47.221	10	
15	4588.20	111.856	13	<b>41.019</b>

# $T=54, m=12$ – Algorithm Opt 1 vs. Algorithm Opt $2 \times 2$

$l$	Algorithm Opt 1		Algorithm Opt $2 \times 2$		$\times$
	Time	No. Sol.	Time	No. Sol.	
21	9.360	12	3.416	2	<b>2.740</b>
22	29.640	12	3.837	2	<b>7.725</b>
23	55.785	17	4.852	4	<b>11.497</b>
24	58.157	17	5.054	4	<b>11.507</b>
25	210.522	20	5.694	6	<b>36.972</b>
26	3428.123	30	20.498	7	<b>167.241</b>
27			21.481	7	
28			88.795	9	
29			510.557	13	

## **T=54, m=12 – CPLEX 11.2 vs. Algorithm Opt $2 \times 2$**

$l$	CPLEX 11.2	Algorithm Opt $2 \times 2$		$\times$
	Time	Time	No. Sol.	
21	14274 3:57:54	3.416	2	<b>4178.57</b>

In[698]:=

