# Automated Solving of Triangle Construction Problems - ongoing work - 

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## Constructions with Straightedge and Compass

- Goal: construct a triangle that meets given constraints
- Widely studied on all education levels
- Main obstacle: combinatorial explosion - huge search space:
- many different construction steps
- plenty of objects that each step could be applied to
- The construction has to be accompanied by a proof that it meets the given specification

Constructions with Straightedge and Compass
Our Approach
Discussion Future Work

Constructions with Straightedge and Compass

## Example Problem

Problem Solution
Existing Approaches
Wernick's Problems

## Example Problem

Problem: Construct a triangle $A B C$ given vertices $A$ and $B$ and the barycenter $G$

## Problem Solution



Solution: Construct the midpoint $M_{c}$ of the segment $A B$; then construct the vertex $C$ such that $M_{c} G: M_{c} C=1 / 3$

## Existing Approaches

- Just a couple of existing approaches, including:
- Gao and Chou (1998)
- Schreck (2001)
- Gulwani et.al (2011)


## Wernick's Problems

- Created in 1982, some variants in the meanwhile
- Task: construct a triangle given three located points selected from the following list:
- $A, B, C$ - vertices
- $I, O$ - incenter and circumcenter
- H, G - orthocenter and barycenter
- $M_{a}, M_{b}, M_{c}$ - the side midpoints
- $H_{a}, H_{b}, H_{c}$ - feet of vertices on the opposite sides
- $T_{a}, T_{b}, T_{c}$ - intersections of the internal angles bisectors with the opposite sides


## Wernick's Problems (2)

139 non-trivial, significantly different, problems; 25 redundant (R) or locus-restricted (L); some solvable (S), some unsolvable (U); 15 still with unknown status

|  | 57. $A, H, I \quad \mathrm{~S}[9]$ | 85. $M_{a}, M_{b}, H_{a} \mathrm{~S}$ | 113. $M_{a}, T_{b}, T_{c}$ |
| :---: | :---: | :---: | :---: |
|  | $\left.1, A, T_{a}, T_{b} \quad \mathrm{~S} \quad 9\right]$ | 86. $M_{a}, M_{b}, H_{c} \mathrm{~S}$ | 114. $M_{a}, T_{b}, I \quad \mathrm{U}[9]$ |
|  | $1 T_{a}, I \quad \mathrm{~L}$ | 87. $M_{a}, M_{b}, H$ S [9] | 115. $G, H_{a}, H_{b}$ U [9] |
| $A, B_{1}, M_{a}$ | $T_{b}, T_{c} \quad \mathrm{~S}$ | 88. $M_{a}, M_{b}, T_{a}$ U [9] | 116. $G, H_{a}, H \quad \mathrm{~S}$ |
|  | $I \quad \mathrm{~S}$ | 89. $M_{a}, M_{b}, T_{c}$ U [9] | 117. $G, H_{a}, T_{a} \quad \mathrm{~S}$ |
| $A, \quad B, \quad H_{C}$ | $M_{b}$ S | 90. $M_{a}, M_{b}, I \quad \mathrm{U}[10]$ | 118. $G, H_{a}, T_{b}$ |
|  | G S | 91. $M_{a}, G, H_{a} \mathrm{~L}$ | 119. G, Ha, I |
|  | $\psi_{a} \quad \mathrm{~L}$ | 92. $M_{a}, G, H_{b}$ S | 120. G, H, Ta U [9] |
| $A, E$ | $\bigcirc \mathrm{S}$ | 93. $M_{0}, G, H \quad \mathrm{~S}$ | 121. $G, H, I \quad \mathrm{U}[9]$ |
|  | S | 94. $M_{a}, G, T_{a} \mathrm{~S}$ | 122. $G, T_{a}, T_{b}$ |
|  | L | 95. $M_{a}, G, T_{b} \cup[9]$ | 123. $G, T_{a}, I$ |
| $A, B_{1}, F_{Q}$ | U [9] | 96. $M_{a}, G, I \quad \mathrm{~S}[9]$ | 124. $H_{a}, H_{b}, H_{c} \mathrm{~S}$ |
|  | d S | 97. $M_{a}, H_{a}, H_{b} \mathrm{~S}$ | 125. $H_{a}, H_{b}, H$ S |
|  | S | 98. $M_{a}, H_{a}, H \quad \mathrm{~L}$ | 126. $H_{a}, H_{b}, T_{a} \mathrm{~S}$ |
|  | R | 99. $M_{a}, H_{a}, T_{a}$ L | 127. $H_{a}, H_{\mathrm{b}}, T_{c}$ |
|  | U [9] | 100. $M_{a}, H_{a}, T_{b}$ U [9] | 128. $H_{a}, H_{b}, I$ |
| $A, E, F$ | U 99 | 101. $M_{a}, H_{a}, I \quad \mathrm{~S}$ | 129. $H_{a}, H, T_{a} \mathrm{~L}$ |
|  | $H_{b}$ U 9 | 102. $M_{c}, H_{b}, H_{c} \mathrm{~L}$ | 130. $H_{a}, H, T_{b}$ U [9] |
|  | , H S | 103. $M_{a}, H_{b}, H \quad \mathrm{~S}$ | 131. $H_{a}, H, I \quad \mathrm{~S}[9]$ |
| $A, B, \Gamma_{a}$ | $H_{a}, T_{a} \quad \mathrm{~S}$ | 104. $M_{a}, H_{b}, T_{a} \mathrm{~S}$ | 132. $H_{a}, T_{a}, T_{b}$ |
|  | $H_{c}, T_{b}$ | 105. $M_{a}, H_{b}, T_{b} \mathrm{~S}$ | 133. $H_{a}, T_{a}, I \quad \mathrm{~S}$ |
|  | $\xrightarrow{\sim} H_{a}, I$ | 106. $M_{a}, H_{b}, T_{c}$ U [9] | 134. $H_{a}, T_{b}, T_{c}$ |
| $\frac{9}{25 .} A, \quad B, \quad T_{C}$ | F. $O, H_{\mathrm{s}} T_{a} \quad \mathrm{U}[9]$ | 107. $M_{a}, H_{b}, I \quad$ U [9] | 135. $H_{a}, T_{b}, I$ |
|  | 80. O, H, I $\quad$ U [9] | 108. $M_{a}, H, T_{a}$ U [9] | 136. $H, T_{a}, T_{b}$ |
|  | 81. $O, T_{a}, T_{b}$ | 109. $M_{a}, H, T_{b}$ U [10] | 137. $H, T_{a}, I$ |
|  | 82. $O, T_{a}, I \quad \mathrm{~S}[9]$ | 110. $M_{a}, H, I \quad \mathrm{U}[10]$ | 138. $T_{a}, T_{b}, T_{c}$ U [11] |
| 27. $A, M_{a}, I \quad \mathrm{~S}$ [9]\||55. $A, H, T_{a} \quad \mathrm{~S}$ | 83. $M_{a}, M_{b}, M_{c} \mathrm{~S}$ | 111. $M_{a}, T_{a}, T_{b}$ U [10] | 139. $T_{a}, T_{b}, I \quad \mathrm{~S}$ |
| 28. $A, M_{b}, M_{c} \mathrm{~S}$ 年 56. $A, H, T_{b}$ U [9] | 84. $M_{a}, M_{b}, G \mathrm{~S}$ | 112. $M_{a}, T_{a}, I \quad \mathrm{~S}$ |  |

## Basic Approach (1)

- Following careful analysis of all solutions
- Constructions consist of high-level construction steps (for example: if barycenter $G$ and circumcenter $O$ are known, then the orthocenter $H$ can be constructed)
- Simple forward chaining mechanism for search procedure
- Points - only basic objects; lines and circles defined as functions of their points
- Implemented in Prolog


## Basic Approach (2)

- Around 70 general rules used
- Example: if two triangle vertices are given, then the side bisector can be constructed
- For symmetric predicates, no redundant facts are derived
- Solves 60 examples from Wernick's list, each in less than 1s and with the maximal search depth 9
- But... there are too many rules! (it is not problem to search over them, but to invent them)


## Separation of concepts definitions, lemmas, construction steps (1)

Motivating example: Construct the midpoint $M_{c}$ of $A B$ and then construct $C$ such that $M_{c} G: M_{c} C=1: 3$ uses the facts:

- $M_{c}$ is the side midpoint of $A B$
- $G$ is the barycenter of $A B C$
- it holds that $M_{c} G=1 / 3 M_{c} C$
- given points $X$ and $Y$, it is possible to construct the midpoint of the segment $X Y$
- given points $X$ and $Y$, it is possible to construct a point $Z$, such that: $X Y: X Z=1: 3$


## Separation of concepts definitions, lemmas, construction steps (2)

Motivating example: Construct the midpoint $M_{c}$ of $A B$ and then construct $C$ such that $M_{c} G: M_{c} C=1: 3$ uses the facts:

- $M_{c}$ is the side midpoint of $A B$ (definition of $M_{c}$ )
- $G$ is the barycenter of $A B C$ (definition of $G$ )
- it holds that $M_{c} G=1 / 3 M_{c} C$ (lemma)
- given points $X$ and $Y$, it is possible to construct the midpoint of the segment $X Y$ (construction primitive)
- given points $X$ and $Y$, it is possible to construct a point $Z$, such that: $X Y: X Z=1: 3$ (construction primitive)


## Advanced Approach

- Task: Derive high-level (instantiated) construction steps from the set of definitions, lemmas and construction primitives
- From:
- it holds that $M_{c} G=1 / 3 M_{c} C$ (lemma)
- given points $X$ and $Y$, it is possible to construct a point $Z$, such that: $X Y: X Z=1: r$ (construction primitive) we can derive:
- given $M_{c}$ and $G$, it is possible to construct $C$


## Rule derivation

- Limit instantiations of definitions/lemmas
- So far, half of the rules of the basic system are derived from:
- around 15 definitions (including Wernick's notation)
- around 10 lemmas
- only 2 suitable construction primitives
- Deriving rules is performed once, in preprocessing phase (takes approx. 20s)


## Discussion

- Objection: the approach is problem-tailored!
- Answer: no system can invent all needed lemmas, so other systems are too problem-tailored
- Objection: how can the approach be used for other families of problems?
- Answer: in analogy with this family (the knowledge may overlap partly)
- Objection: ...then, it might become inefficient?
- Answer: It could automatically choose over domains


## Future Work

- Complete the process of automated deriving of rules
- Automated generation of constructions and figures in GCLC (along with a construction description in ${ }^{A} T_{E X}$ )
- Proving (in GCLC) that the constructions meet specifications, using automated theorem provers
- Proving (in coherent logic, by ArgoCLP prover) that constructed points indeed exist (under some conditions)

