Alphageometry in Practice

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Language

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"After running this process on 100,000 CPU workers for 72h, we obtained roughly 500 million synthetic proof examples. We reformat the proof statements to their canonical form (for example, sorting arguments of individual terms and sorting terms within the same proof step, etc.) to avoid shallow deduplication against itself and against the test set. At the end, we obtain 100 million unique theorem–proof examples. A total of 9 million examples involves at least one auxiliary construction. We find no IMO-AG-30 problems in the synthetic data. On the set of geometry problems collected in JGEX17, which consists mainly of problems with moderate difficulty and well-known theorems, we find nearly 20 problems in the synthetic data. This suggests that the training data covered a fair amount of common knowledge in geometry, but the space of more sophisticated theorems is still much larger."

What is the starting point?

https://github.com/google-deepmind/alphageometry/blob/main/rules.txt

Rules

1. perp A B C D, perp C D E F, ncoll A B E => para A B E F 2. cong O A O B, cong O B O C, cong O C O D => cyclic A B C D 3. eqangle A B P Q C D P Q => para A B C D 4. cyclic A B P Q => eqangle P A P B Q A Q B 5. eqangle6 P A P B Q A Q B, ncoll P Q A B => cyclic A B P Q 6. cyclic A B C P Q R, eqangle C A C B R P R Q => cong A B P Q 7. midp E A B, midp F A C => para E F B C 8. para A B C D, coll O A C, coll O B D => eqratio3 A B C D O O 9. perp A B C D, perp E F G H, npara A B E F => eqangle A B E F C D G H 10.eqangle a b c d m n p q, eqangle c d e f p q r u => eqangle a b e f m n r u 11.eqratio a b c d m n p q, eqratio c d e f p q r u => eqratio a b e f m n r u 12.eqratio6 d b d c a b a c, coll d b c, ncoll a b c => eqangle6 a b a d a d a 13. eqangle6 a b a d a d a c, coll d b c, ncoll a b c => eqratio6 d b d c a b a c c 14.cong O A O B, ncoll O A B => eqangle O A A B A B O B 15.eqangle6 A O A B B A B O, ncoll O A B => cong O A O B

Rules

16 cong $0 A 0 B$, ncoll $0 A B \Rightarrow$ eqangle $0 A A B A B 0 B$ 17 eqangle6 A O A B B A B O, ncoll O A B => cong O A O B 18 circle O A B C, perp O A A X => eqangle A X A B C A C B 19 circle O A B C, eqangle A X A B C A C B => perp O A A X 20 circle 0 A B C, midp M B C => eqangle A B A C 0 B 0 M 21 circle 0 A B C, coll M B C, eqangle A B A C 0 B 0 M => midp M B C 22 perp A B B C, midp M A C => cong A M B M 23 circle O A B C, coll O A C => perp A B B C 24 cyclic A B C D, para A B C D => eqangle A D C D C D C B 25 cyclic A B C D, para A B C D => eqangle A D C D C D C B 26 midp M A B, perp O M A B = $>$ cong O A O B 27 cong A P B P, cong A Q B Q => perp A B P Q 28 cong A P B P, cong A Q B Q, cyclic A B P Q => perp P A A Q 29 midp M A B, midp M C D => para A C B D 30 midp M A B, para A C B D, para A D B C => midp M C D 31 eqratio 0 A A C 0 B B D, coll 0 A C, coll 0 B D, ncoll A B C, sameside A 0 C B 0 D => para A B C D 32.para $A \text{ } B \text{ } A \text{ } C \Rightarrow$ coll $A \text{ } B \text{ } C$

Rules

33. midp M A B, midp N C D => egratio M A A B N C C D 34. eqangle A B P Q C D U V, perp P Q U V => perp A B C D 35. eqratio A B P Q C D U V, cong P Q U V => cong A B C D 36. cong A B P Q, cong B C Q R, cong C A R P, ncoll A B C => contri^{*} A B C P Q R 37. cong A B P Q, cong B C Q R, eqangle6 B A B C Q P Q R, ncoll A B C => contri* A B C P Q R 38. eqangle6 B A B C Q P Q R, eqangle6 C A C B R P R Q, ncoll A B C => simtri A B C P Q R 39. eqangle6 B A B C Q R Q P, eqangle6 C A C B R Q R P, ncoll A B C => simtri2 A B C P Q R 40. eqangle6 B A B C Q P Q R, eqangle6 C A C B R P R Q, ncoll A B C, cong A B P Q => contri A B C P Q R 41. eqangle6 B A B C Q R Q P, eqangle6 C A C B R Q R P, ncoll A B C, cong A B P Q => contri2 A B C PQR 42. eqratio6 B A B C Q P Q R, eqratio6 C A C B R P R Q, ncoll A B C => simtri* A B C P Q R 43. eqratio6 B A B C Q P Q R, eqangle6 B A B C Q P Q R, ncoll A B C => simtri* A B C P Q R 44. eqratio6 B A B C Q P Q R, eqratio6 C A C B R P R Q, ncoll A B C, cong A B P Q => contri* A B C P Q R 45. para a b c d, coll m a d, coll n b c, eqratio6 m a m d n b n c, sameside m a d n b c => para m n a b

46. para a b c d, coll m a d, coll n b c, para m n a b => eqratio6 m a m d n b n c

How Alphageometry Solves Problems

•**Deductive Reasoning (DD)**: The system starts by applying geometric rules to draw new conclusions. For example, it might deduce that two lines, AB and CD, are parallel.

•**Algebraic Reasoning (AR)**: The information deduced by DD, such as "AB is parallel to CD," is then translated into algebraic form. In this case, it means that the slopes of lines AB and CD will be set as equal in a matrix of equations managed by AR.

•**Algebraic Processing**: The system uses methods like Gaussian elimination to manipulate these equations and find new relationships between variables. These new relationships are then passed back to the deductive part (DD).

•**Repetitive Process**: This process of passing conclusions back and forth between DD and AR continues until the system can't make any more new deductions (when the "deduction closure" stops expanding).

Our paper shows that language models can learn to come up with auxiliary constructions from synthetic data, in which problem statements and auxiliary constructions are randomly generated together and then separated using the traceback algorithm to identify the dependency difference.

Additional constructs used in solving IMO tasks

Additional constructs

- Construct point X as the midpoint of segment AB
- Construct point X as the mirror of A through B.
- Construct point X as the intersection of circles (O, A) and (O1, A).
- Construct point X as the circumcenter of triangle ABC.
- Construct point X as the orthocenter of triangle ABC.
- Construct point X such that XA = XB.
- Construct point X such that $AB = BX$ and AX is perpendicular to BC.
- Construct point X such that $AX = BX$ and $ACAX = CXCA$.
- Construct point X such that $AB = CX$ and $AX = CB$.
- Construct point X such that AB is parallel to CX and AX is parallel to CD.
- Construct point X as the foot of A on line BC.

• Construct point X such that AB = BX and AX is perpendicular to BC.

Construct point X such that AB is parallel to CX and AX is parallel to CD

Sample tasks: IMO 2000 P1

Original:

Two circles G1 and G2 intersect at two points M and N. Let AB be the line tangent to these circles at A and B, respectively, so that M lies closer to AB than N. Let CD be the line parallel to AB and passing through the point M, with C on G1 and D on G2. Lines AC and BD meet at E; lines AN and CD meet at P; lines BN and CD meet at Q. Show that EP = EQ.

Sample tasks: IMO 2000 P1

Original:

Two circles G1 and G2 intersect at two points M and N. Let AB be the line tangent to these circles at A and B, respectively, so that M lies closer to AB than N. Let CD be the line parallel to AB and passing through the point M, with C on G1 and D on G2. Lines AC and BD meet at E; lines AN and CD meet at P; lines BN and CD meet at Q. Show that $EP = EQ$.

Translated (in paper):

Let A and B be any two distinct points. Define point G1 such that AB is perpendicular to AG1. Define point G2 such that AB is perpendicular to BG2. Define point M as the intersection of circles (G1, A) and (G2, B). Define point N as the intersection of circles (G1, A) and (G2, B). Define point C on circle (G1, A) such that AB is parallel to CM. Define point D on circle (G2, B) such that AB is parallel to DM. Define point E as the intersection of lines AC and BD. Define point P as the intersection of lines AN and CD. Define point Q as the intersection of lines BN and CD. Prove that $EP = EQ$

Sample tasks: IMO 2000 P1

Alphageometry:

a b = segment a b; $g1 =$ on_tline g1 a a b; $g2$ = on_tline $g2$ b b a; $m =$ on_circle m g1 a, on_circle m g2 b; n = on_circle n g1 a, on_circle n g2 b; c = on_pline c m a b, on_circle c g1 a; $d =$ on pline d m a b, on circle d g2 b; e = on_line e a c, on_line e b d; p = on_line p a n, on_line p c d; $q =$ on_line q b n, on_line q c d ? cong e p e q

Original (in paper): Sample tasks: IMO 2004 P5A

In a convex quadrilateral ABCD, the diagonal BD bisects neither the angle ABC nor the angle CDA. The point P lies inside ABCD and satisfies \leq PBC = \leq DBA and \leq PDC = \leq BDA. Prove that AP=CP given ABCD is a cyclic quadrilateral.

Original: Prove that ABCD is a cyclic quadrilateral if and only if AP = CP.

Translated:

Let ABC be a triangle. Define point O as the circumcenter of triangle CBA. Let D be any point on circle (O, A). Define point P such that <ABD = <PBC and <ADB = <PDC. Prove that AP = CP

Alphageometry:

a $b c =$ triangle a $b c$;

o = circle o a b c;

d = on_circle d o a;

p = on_aline p b c a b d, on_aline p d c a d b

? cong a p c p

Figure by Alphageometry

Sample tasks: IMO 2004 P5A

Step 1. AO = BO, AO = DO and BO = $CO \Rightarrow A$, B, C, D are cyclic.

Step 2. A, B, C, D are cyclic \Rightarrow <BAD = <BCD and <BAC = <BDC.

Step 3. AO = BO, AO = DO and BO = $CO \Rightarrow CO = DO$.

Step 4. $CO = DO \Rightarrow \text{CDO} = \text{COC}$.

Step 5. $BO = CO \Rightarrow .$

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Step 6. <BAD = <BCD, <ABD = <PBC, <BCO = <OBC, <ADB = 
\leqPDC and \leqCDO = \leqOCD \Rightarrow by angle chasing: \leqBOD = \leqBPD.
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Step 7. \leqBOD = \leqBPD \Rightarrow B, D, O, P are cyclic.
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Step 8. B, D, O, P are cyclic \Rightarrow <BDP = <BOP.
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Step 9. AO = BO and BO = CO \Rightarrow AO = CO.
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Step 10. AO = CO \Rightarrow <ACO = <OAC.
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Step 11. AO = BO \Rightarrow <ABO = <OAB.
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Step 12. <BAC = <BDC, <BAD = <BCD, <ABO = <OAB, <ACO =
\leqOAC, \leqBCO = \leqOBC, \leqADB = \leqPDC and \leqBDP = \leqBOP \Rightarrow by
angle chasing: OP is the bisector of <AOC.
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Step 13. AO = CO and OP is the bisector of \leq AOC \Rightarrow AP = CP
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https://www.cut-the-knot.org/pythagoras/IMO2004-5.shtml

My tasks

1. a b c = triangle a b c; d = on_pline d c b a, on_pline d b c a ? equangle a b c c d a

2. a b = segment; c = on_tline c a b a; d = on_tline d b b a, on_tline d c b a ? eqdistance d a c b

3. a b c = triangle a b c; d = midpoint d a b; e = foot e b c d; f = foot f a c d ? eqdistance a f e b

4. c b a = triangle c b a; $d =$ parallelogram c b a d; o = circumcenter o a d b; $l =$ intersection_lc l c o d; $k =$ intersection_lc k b o c; n = mirror n a o ? eqdistance n k c n

https://www.tandfonline.com/eprint/MKJWRQFEGGEA6HN3RRWB/full?target=10.1080/0020739 X.2024.2377724#d1e288

My task

a b c = triangle a b c; d = midpoint d c b; e = midpoint e c a; $f =$ on_bline f a c, on_bline f c b; g = intersection_ll g a d b e; $h =$ foot h b a c; i = foot i a b c; j = intersection_ll j a i b h ? coll j g f

Citations from the article

This auxiliary construction can be found quickly with the knowledge of Reim's theorem, which is not included in the deduction rule list used by the symbolic engine during synthetic data generation. Including such high-level theorems into the synthetic data generation can greatly improve the coverage of synthetic data and thus improve auxiliary construction capability. Further, higher-level steps using Reim's theorem also cut down the current proof length by a factor of 3.

AlphaGeometry constructs point K to materialize this axis, whereas humans simply use the existing point R for the same purpose. This is a case in which proof pruning itself cannot remove K and a sign of similar redundancy in our synthetic data.

This human proof uses four auxiliary constructions (diameters of circles W1 and W2) and highlevel theorems such as the Pitot theorem and the notion of homothety. These high-level concepts are not available to our current version of the symbolic deduction engine both during synthetic data generation and proof search. Again, this suggests that enhancing the symbolic engine with more powerful tools that IMO contestants are trained to use can improve both the synthetic data and the test-time performance of AlphaGeometry.

Pitot theorem

The Pitot theorem in geometry states that in a tangential quadrilateral the two pairs of opposite sides have the same total length.

https://en.wikipedia.org/wiki/Pitot_theorem

(Reim's Theorem). Choose points A, B, X, Y on circle ω1 and let C and D be points on AX and BY . Then AB || CD if X, Y, C, D are concyclic.

AlphaGeometry excels at solving problems involving cyclic quadrilaterals, with cyclic points appearing in 24 out of 25 solved tasks.

This suggests that the system is particularly strong when working with such geometric structures. However, this raises the possibility that the tasks were selected to play to the system's strengths. It's conceivable that a different set of problems, without reliance on cyclic figures, could be more challenging for AlphaGeometry but potentially solvable by other provers.